

# On Hitting All Maximum Cliques with an Independent Set

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**Abstract:** We prove that every graph  $G$  for which  $\omega(G) \geq \frac{3}{4}(\Delta(G) + 1)$  has an independent set  $I$  such that  $\omega(G - I) < \omega(G)$ . It follows that a minimum counterexample  $G$  to Reed's conjecture satisfies  $\omega(G) < \frac{3}{4}(\Delta(G) + 1)$  and hence also  $\chi(G) > \lceil \frac{7}{6}\omega(G) \rceil$ . This also applies to restrictions of Reed's conjecture to hereditary graph classes, and in particular generalizes and simplifies King, Reed and Vetta's proof of Reed's conjecture for line graphs. © 2010 Wiley Periodicals, Inc. *J Graph Theory* 66: 32–37, 2010

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## 1. INTRODUCTION

We prove the following general lemma and apply it to Reed's conjecture.

**The Main Lemma.** *If  $G$  is a graph with  $\omega(G) \geq \frac{3}{4}(\Delta(G) + 1)$ , then  $G$  has an independent set  $I$  such that  $\omega(G - I) < \omega(G)$ .*

In [9], Reed conjectured the following upper bound on the chromatic number.

**Reed's Conjecture.** *For every graph  $G$  we have  $\chi(G) \leq \lceil (\omega(G) + \Delta(G) + 1)/2 \rceil$ .*

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**Observation.** If we could always find an independent set whose removal decreased both  $\omega$  and  $\Delta$ , then the conjecture would follow by simple induction since we can give the independent set a single color and use at most  $\lceil(\omega(G)+\Delta(G)+1)/2\rceil - 1$  colors on what remains. Expanding the independent set given by The Main Lemma to a maximal one shows that this sort of argument goes through when  $\omega \geq \frac{3}{4}(\Delta+1)$ . Thus, a minimum counterexample to Reed’s conjecture satisfies  $\omega < \frac{3}{4}(\Delta+1)$  and hence also  $\chi > \lceil \frac{7}{6}\omega \rceil$ .

Reed’s conjectured upper bound was proved for line graphs of multigraphs by King, Reed and Vetta in [4], for quasi-line graphs by King and Reed in [5], and recently King and Reed proved it for all claw-free graphs (see King’s thesis [6]). The line graphs of multigraphs result follows from the following theorem.

**Theorem D.** *If  $G$  is a graph with  $\chi(G) \leq \lceil \frac{7}{6}\omega(G) \rceil$  and for every proper induced subgraph  $H$  of  $G$  we have  $\chi(H) \leq \lceil(\omega(H)+\Delta(H)+1)/2\rceil$ , then we also have  $\chi(G) \leq \lceil(\omega(G)+\Delta(G)+1)/2\rceil$ .*

King, Reed and Vetta’s upper bound for line graphs of multigraphs follows immediately from Theorem D, a bound of Caprara and Rizzi (see [1]) and the bound of Molloy and Reed on fractional colorings (see [8, Chapter 21, Section 3]). We write  $\chi^*(G)$  for the fractional chromatic number of a graph  $G$ .

**Caprara and Rizzi.** *Let  $H$  be a multigraph and  $G=L(H)$ . Then*

$$\chi(G) \leq \max\{\lfloor 1.1\Delta(H)+0.7 \rfloor, \lceil \chi^*(G) \rceil\}$$

**Fractional Version.** *For every graph  $G$  we have  $\chi^*(G) \leq (\omega(G)+\Delta(G)+1)/2$ .*

**King, Reed and Vetta.** *If  $G$  is the line graph of a multigraph, then  $\chi(G) \leq \lceil(\omega(G)+\Delta(G)+1)/2\rceil$ .*

To prove this result, they consider a minimum counterexample  $G=L(H)$ . This must have  $\Delta(G) < \frac{3}{2}\Delta(H) - 1$ ; otherwise, Caprara and Rizzi’s result implies Reed’s bound. When  $\Delta(G) < \frac{3}{2}\Delta(H) - 1$ , they use the structure of line graphs to prove that  $G$  contains an independent set intersecting all maximum cliques. However, this second step can easily be replaced with an application of Theorem D, since  $\lfloor 1.1\omega(G)+0.7 \rfloor \leq \lceil \frac{7}{6}\omega(G) \rceil$ . This example demonstrates the usefulness of Theorem D as a general tool, requiring no specific structural analysis to find the independent set. On the other hand, the structural analysis has the benefit of leading to a polynomial time algorithm for actually finding the desired independent set in a supposedly minimum counterexample (see [4]).

## 2. PROOF OF THE MAIN LEMMA

We need three lemmas. The first is due to Hajnal (see [2]).

**Lemma 1.** *Let  $G$  be a graph and  $\mathcal{Q}$  a collection of maximum cliques in  $G$ . Then*

$$|\cap \mathcal{Q}| \geq 2\omega(G) - |\cup \mathcal{Q}|.$$

**Proof.** Assume (to reach a contradiction) that the lemma is false and let  $\mathcal{Q}$  be a counterexample with  $|\mathcal{Q}|$  minimal. Put  $r=|\mathcal{Q}|$  and  $\mathcal{Q}=\{Q_1, \dots, Q_r\}$ . Consider the set  $W=(Q_1 \cap \bigcup_{i=2}^r Q_i) \cup \bigcap_{i=2}^r Q_i$ . Plainly,  $W$  is a clique. Thus

$$\begin{aligned} \omega(G) &\geq |W| \\ &= \left| (Q_1 \cap \bigcup_{i=2}^r Q_i) \cup \bigcap_{i=2}^r Q_i \right| \\ &= \left| Q_1 \cap \bigcup_{i=2}^r Q_i \right| + \left| \bigcap_{i=2}^r Q_i \right| - \left| \bigcap_{i=1}^r Q_i \cap \bigcup_{i=2}^r Q_i \right| \\ &= |Q_1| + \left| \bigcup_{i=2}^r Q_i \right| - \left| \bigcup_{i=1}^r Q_i \right| + \left| \bigcap_{i=2}^r Q_i \right| - \left| \bigcap_{i=1}^r Q_i \right| \\ &= \omega(G) + \left| \bigcup_{i=2}^r Q_i \right| + \left| \bigcap_{i=2}^r Q_i \right| - \left| \bigcup_{i=1}^r Q_i \right| - \left| \bigcap_{i=1}^r Q_i \right| \\ &\geq \omega(G) + 2\omega(G) - \left| \bigcup_{i=1}^r Q_i \right| - \left| \bigcap_{i=1}^r Q_i \right|, \end{aligned}$$

where the final step follows by the minimality of  $|\mathcal{Q}|$ . Thus  $|\bigcap_{i=1}^r Q_i| \geq 2\omega(G) - |\bigcup_{i=1}^r Q_i|$  giving a contradiction. ■

The second lemma we need is an improvement of Hajnal’s result for graphs satisfying  $\omega > \frac{2}{3}(\Delta + 1)$  due to Kostochka (see [7]). We reproduce Kostochka’s proof here to serve as an English translation.

**Clique Graph.** Let  $G$  be a graph and  $\mathcal{Q}$  the collection of all maximum cliques in  $G$ . The clique graph of  $G$  is the graph with vertex set  $\mathcal{Q}$  and an edge between  $Q_1 \neq Q_2 \in \mathcal{Q}$  if and only if  $Q_1$  and  $Q_2$  intersect. Let  $\mathcal{C}(G)$  be the components of the clique graph of  $G$ .

**Lemma 2.** Let  $G$  be a graph with  $\omega(G) > \frac{2}{3}(\Delta(G) + 1)$ . Then for every  $C \in \mathcal{C}(G)$  we have

$$|\cap V(C)| \geq 2\omega(G) - (\Delta(G) + 1).$$

**Proof.** Assume (to reach a contradiction) that the lemma is false and let  $G$  be a counterexample with the minimum number of vertices. Then there is some component  $C \in \mathcal{C}(G)$  with  $|\cap V(C)| < 2\omega(G) - (\Delta(G) + 1)$ . By minimality of  $G$ , we have  $G = \bigcup V(C)$ . Put  $D = \cap V(C)$ . Note that if  $|D| \geq 1$ , then  $|G| \leq \Delta(G) + 1$  and the result follows from Lemma 1. We can therefore assume that  $|D| = 0$  and hence  $|V(C)| \geq 3$ .

By Lemma 1 we have  $|G| \geq 2\omega(G)$ . Put  $V(C) = \{Q_1, \dots, Q_r\}$ . Take  $x \in V(G)$  that is in the minimum number of the  $Q_i$ . Without loss of generality, say  $x \in Q_i$  for  $1 \leq i \leq t$  for some  $t \geq 1$ . Consider the set

$$A = \bigcap_{i=1}^t Q_i - \bigcup_{i=t+1}^r Q_i.$$

If  $y \in G - A$ , then  $y \notin \bigcap_{i=1}^t Q_i$  or  $y \in \bigcup_{i=t+1}^r Q_i$ . In the former case, we must have  $y \in \bigcup_{i=t+1}^r Q_i$  for otherwise  $y$  would be in fewer than  $t$  of the  $Q_i$  contradicting the minimality of  $x$ . Thus  $G - A \subseteq \bigcup_{i=t+1}^r Q_i$ . Hence  $G - A = \bigcup_{i=t+1}^r Q_i$ .

To apply the minimality of  $G$  to  $G - A$ , all we need to show is that  $G - A$  has a single clique component. Clearly, this will follow if we show that the clique graph of  $G$  is complete. Since the clique graph is connected, it will be enough to show that it is transitive, i.e. contains no induced path on three vertices. So, let  $Q_1, Q_2, Q_3$  be distinct maximum cliques and assume that  $Q_1 \cap Q_2 \neq \emptyset$  and  $Q_2 \cap Q_3 \neq \emptyset$ . Then  $|Q_1 \cap Q_2| = |Q_1| + |Q_2| - |Q_1 \cup Q_2| \geq 2\omega(G) - (\Delta(G) + 1)$ . Hence

$$\begin{aligned} |Q_1 \cap Q_3| &\geq |Q_1 \cap Q_2 \cap Q_3| \\ &\geq |Q_1 \cap Q_2| - (|Q_2| - |Q_2 \cap Q_3|) \\ &\geq 2\omega(G) - (\Delta(G) + 1) - (\omega(G) - (2\omega(G) - (\Delta(G) + 1))) \\ &= 3\omega(G) - 2(\Delta(G) + 1) > 0. \end{aligned}$$

Thus  $Q_1 \cap Q_3 \neq \emptyset$ , showing that the clique graph of  $G$  is transitive.

So we may apply minimality of  $G$  to conclude that  $|\bigcap_{i=t+1}^r Q_i| \geq 2\omega(G) - (\Delta(G) + 1)$ . In particular,  $|G - A| \leq \Delta(G) + 1$ . Since  $A \subseteq Q_1 - Q_r$  and  $Q_1 \cap Q_r \neq \emptyset$ , we have  $|A| \leq \omega(G) - (2\omega(G) - (\Delta(G) + 1)) = \Delta(G) + 1 - \omega(G)$ . But then

$$\begin{aligned} |G| &= |A| + |G - A| \\ &\leq \Delta(G) + 1 - \omega(G) + \Delta(G) + 1 \\ &= 2(\Delta(G) + 1) - \omega(G) \\ &< 2\omega(G). \end{aligned}$$

This contradicts the fact that  $|G| \geq 2\omega(G)$ . ■

Kostochka gives the example of  $C_5$  with each vertex blown up to a  $k$ -clique to show that the  $\omega > \frac{2}{3}(\Delta + 1)$  condition in Lemma 2 is best possible.

The third lemma we need is a result of Haxell (see [3]) on independent transversals.

**Lemma 3.** *Let  $k$  be a positive integer, let  $H$  be a graph of maximum degree at most  $k$  and let  $V(H) = V_1 \cup \dots \cup V_n$  be a partition of the vertex set of  $H$ . Suppose that  $|V_i| \geq 2k$  for each  $i$ . Then  $H$  has an independent set  $\{v_1, \dots, v_n\}$  where  $v_i \in V_i$  for each  $i$ .*

**Proof of The Main Lemma.** Let  $G$  be a graph satisfying  $\omega \geq \frac{3}{4}(\Delta + 1)$ . Put  $\mathcal{C}(G) = \{C_1, \dots, C_r\}$ . By Lemma 2, the mutual intersection  $F_i$  of the maximum cliques in  $C_i$  satisfies  $|F_i| \geq 2\omega(G) - (\Delta(G) + 1)$  for each  $i$ . Since every vertex  $v \in F_i$  is in a maximum clique in  $\bigcup V(C_i)$ ,  $v$  is adjacent to at most  $\Delta(G) + 1 - \omega(G) \leq \frac{1}{4}(\Delta(G) + 1)$  vertices outside of  $\bigcup V(C_i)$ .

Let  $H$  be the graph with  $V(H) = \bigcup_i V(F_i)$  and an edge between  $v, w \in V(H)$  if and only if  $vw \in E(G)$  and  $v$  and  $w$  are in different clique components in  $G$ . Then, by the above,  $\Delta(H) \leq \Delta(G) + 1 - \omega(G)$ .

Consider the partition  $\{F_i\}_i$  of  $V(H)$ . We have

$$\begin{aligned} |F_i| &\geq 2\omega(G) - (\Delta(G) + 1) \\ &\geq 2\frac{3}{4}(\Delta(G) + 1) - (\Delta(G) + 1) \\ &= \frac{1}{2}(\Delta(G) + 1) \\ &\geq 2\Delta(H). \end{aligned}$$

Thus, by Lemma 3,  $H$  has an independent set  $I = \{v_1, \dots, v_n\}$  where  $v_i \in F_i$  for each  $i$ . Since  $F_i$  is contained in all the maximum cliques in  $C_i$ , we have  $\omega(G - I) < \omega(G)$ . ■

### 3. PROOF OF THEOREM D

Theorem D is an easy consequence of The Main Lemma.

**Proof of Theorem D.** Assume (to reach a contradiction) that the theorem is false and let  $G$  be a counterexample with the minimum number of vertices. First assume that  $\omega(G) \geq \frac{3}{4}(\Delta(G) + 1)$ . Then by The Main Lemma we have an independent set  $I$  with  $\omega(G - I) < \omega(G)$ . Plainly, we may assume that  $I$  is maximal (and hence  $\Delta(G - I) < \Delta(G)$ ). Put  $H = G - I$ . Then, by minimality of  $G$ , we have

$$\begin{aligned} \chi(G) &\leq 1 + \chi(H) \\ &\leq 1 + \left\lceil \frac{\omega(H) + \Delta(H) + 1}{2} \right\rceil \\ &\leq 1 + \left\lceil \frac{\omega(G) - 1 + \Delta(G) - 1 + 1}{2} \right\rceil \\ &\leq \left\lceil \frac{\omega(G) + \Delta(G) + 1}{2} \right\rceil. \end{aligned}$$

This is a contradiction; hence, we must have  $\omega(G) < \frac{3}{4}(\Delta(G) + 1)$ . But then

$$\begin{aligned} \left\lceil \frac{7}{6}\omega(G) \right\rceil &\geq \chi(G) \\ &> \left\lceil \frac{\omega(G) + \Delta(G) + 1}{2} \right\rceil \\ &\geq \left\lceil \frac{\omega(G) + \frac{4}{3}\omega(G)}{2} \right\rceil \\ &= \left\lceil \frac{7}{6}\omega(G) \right\rceil. \end{aligned}$$

This final contradiction completes the proof. ■

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